

# Quantum Circuit Completeness: Extensions and Simplifications

32nd EACSL Annual Conference on Computer Science Logic 2024 (CSL'24)

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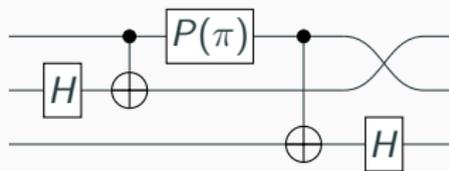
Alexandre Clément\*, Noé Delorme<sup>‡</sup>, Simon Perdrix<sup>‡</sup>, and Renaud Vilmart\*

\* Université Paris-Saclay, ENS Paris-Saclay, CNRS, Inria, LMF, 91190, Gif-sur-Yvette, France

<sup>‡</sup> Université de Lorraine, CNRS, Inria, LORIA, F-54000 Nancy, France

# What is it all about?

Quantum circuits are a rigorous graphical representation of quantum algorithms.

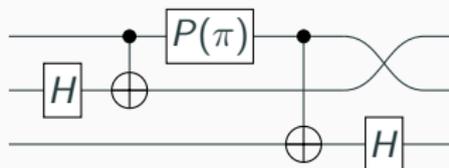


Just like boolean circuits are a rigorous graphical representation of classical algorithms.

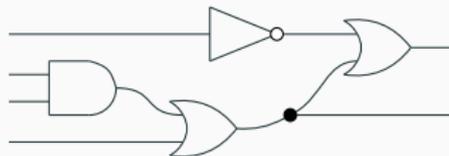


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# Quantum circuits as a graphical language

Quantum circuits are generated by



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to form new circuits.

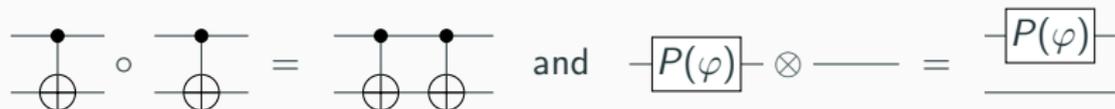


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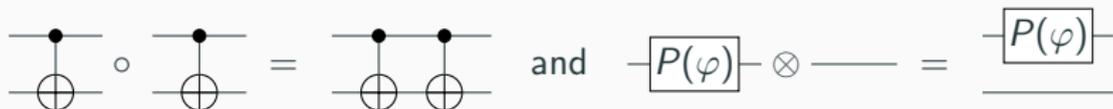


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# Quantum circuits as a graphical language

Formally, graphical languages are defined within the [prop formalism](#) with some deformation rules such that

$$\boxed{P(\varphi)} \circ \text{---} = \boxed{P(\varphi)} \quad \text{or} \quad \text{---} \times \text{---} = \text{---}$$

This framework captures the intuitive behaviour of wires by ensuring that circuits are defined “[up to deformation](#)”.

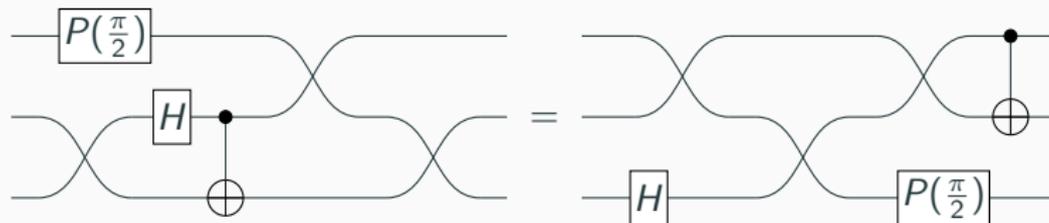


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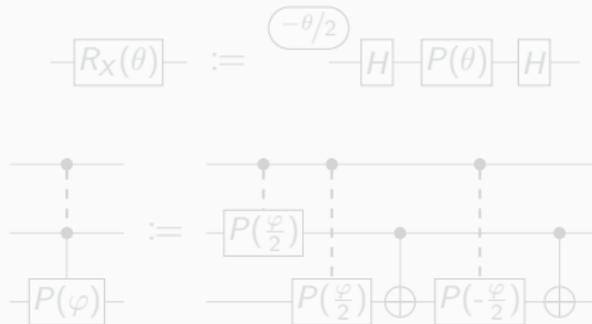


## Other usual gates as shortcut notations

There are only four different kinds of generators



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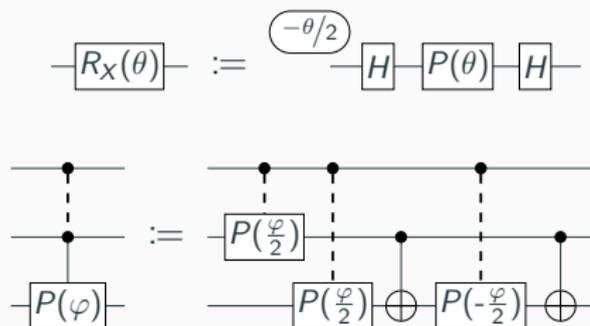


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# Standard interpretation of quantum circuits

## Interpretation

$$[[C_2 \circ C_1]] = [[C_2]] \circ [[C_1]]$$

$$[[C_1 \otimes C_2]] = [[C_1]] \otimes [[C_2]]$$

$$[[\text{I}]] = (1)$$

$$[[\text{Phase}(\varphi)]] = (e^{i\varphi})$$

$$[[\text{CNOT}]] = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

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# Motivations

Distinct circuits can have the same interpretation.

$$\left[ \begin{array}{c} \text{---} P\left(\frac{\pi}{2}\right) \text{---} \bullet \text{---} \text{---} \bullet \text{---} \\ \text{---} P\left(\frac{\pi}{2}\right) \text{---} \oplus \text{---} P\left(-\frac{\pi}{2}\right) \text{---} \oplus \text{---} \end{array} \right] = \left[ \begin{array}{c} \text{---} \bullet \text{---} \\ \text{---} H \text{---} \oplus \text{---} H \text{---} \end{array} \right] = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}$$

Given a quantum algorithm, which circuit is the best?

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- Resource optimisation (for instance the number of gates).
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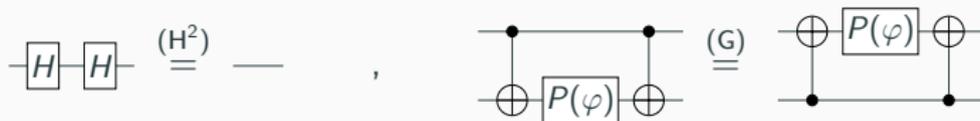
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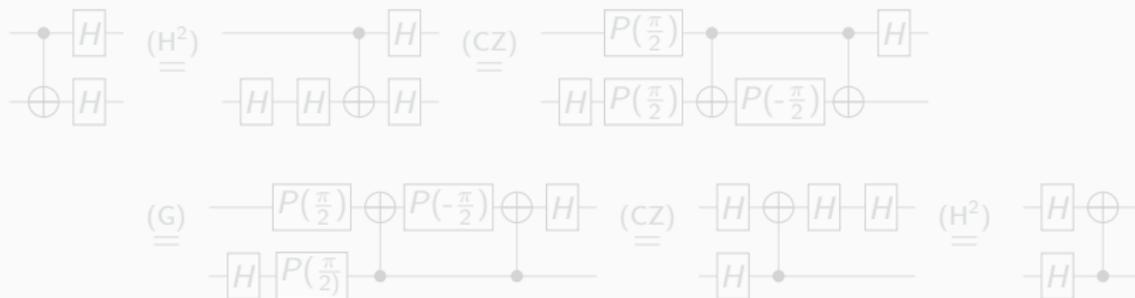
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# Using equations to transform circuits

We can use simple axioms such that,

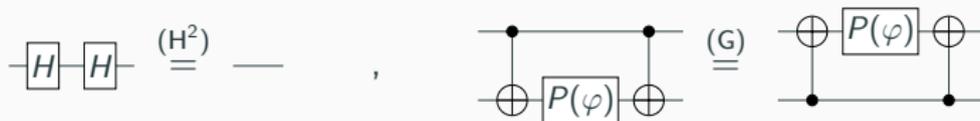


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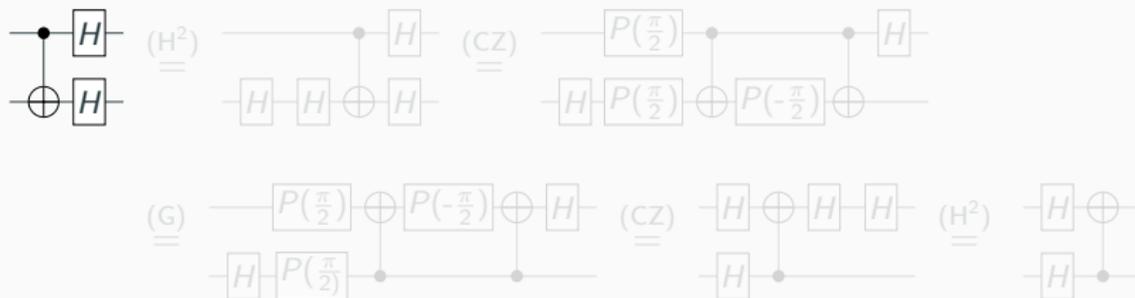


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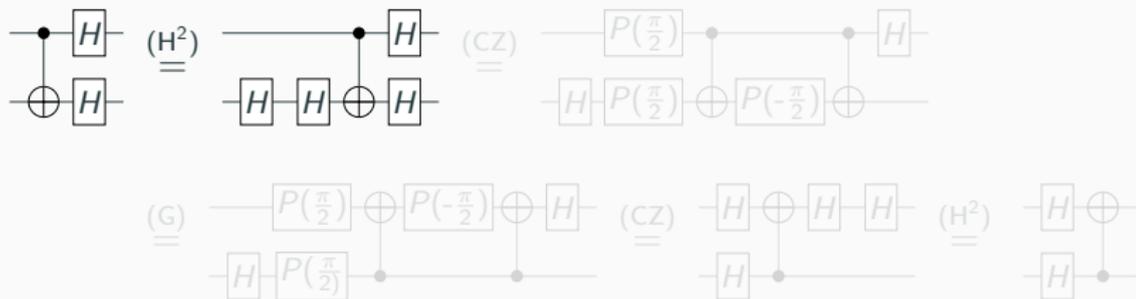


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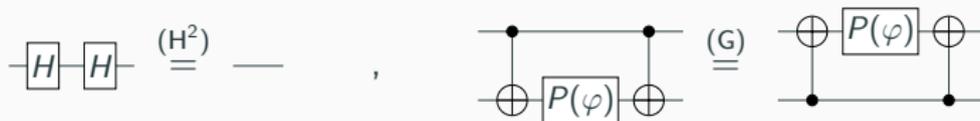


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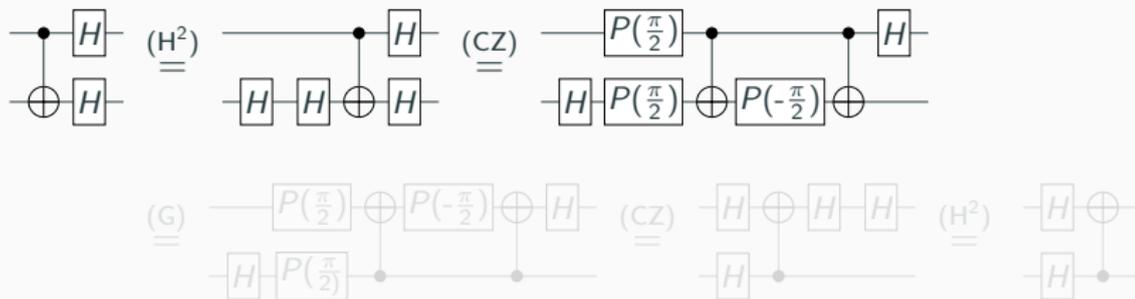


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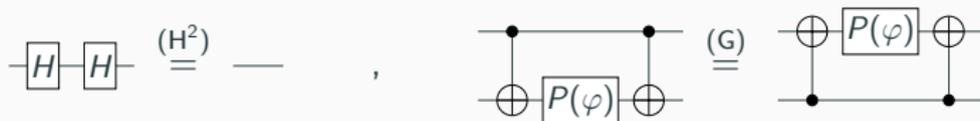


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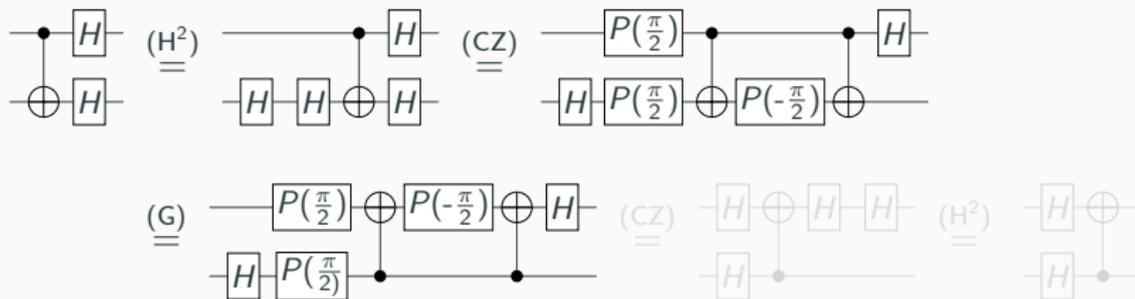


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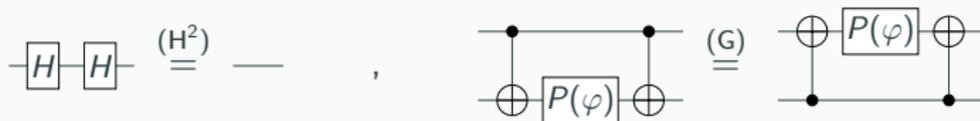


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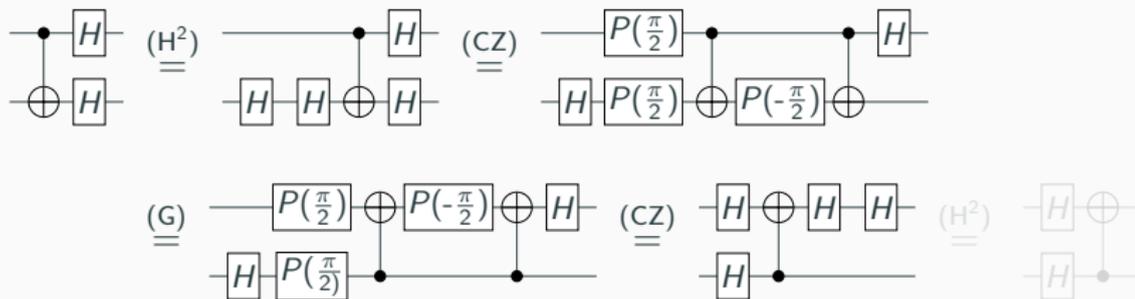


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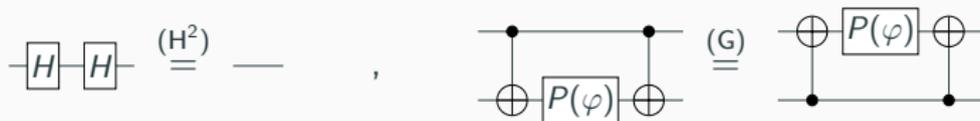


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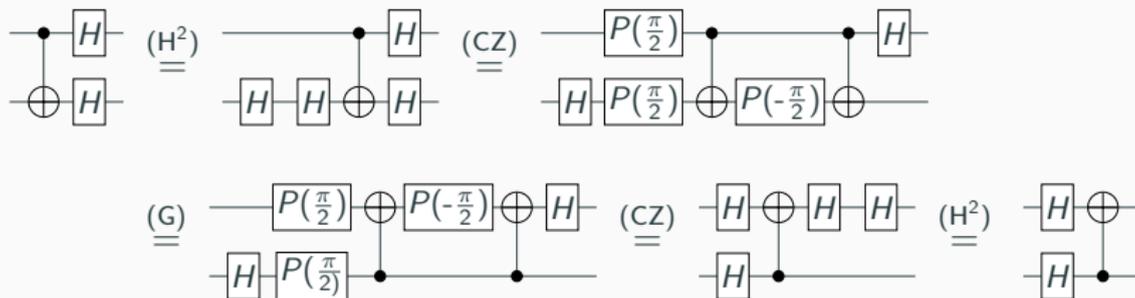


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# Complete and sound equational theory

Is there an **equational theory** (i.e. a set of axioms)  $\Gamma$  from which we can derive any true equation and only true equations?

## Soundness

Any derivable equation is true.

$$\forall C_1, C_2 \quad : \quad \Gamma \vdash C_1 = C_2 \implies \llbracket C_1 \rrbracket = \llbracket C_2 \rrbracket$$

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Previous work [Clément, Heurtel, Mansfield, Perdrix, Valiron LICS'23]:

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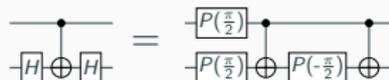
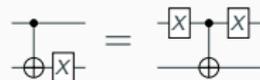


First contribution

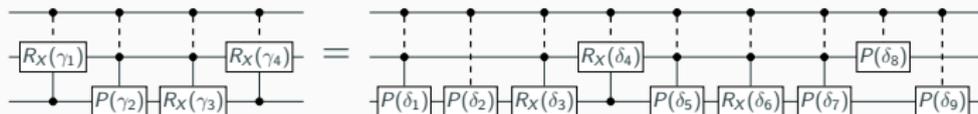
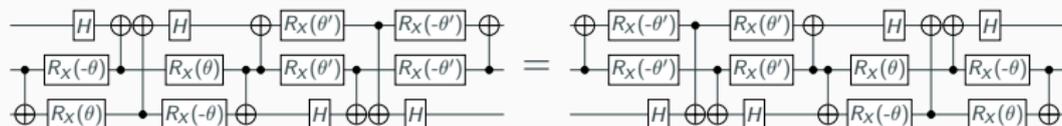
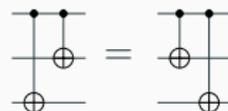
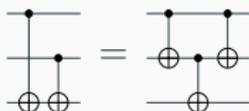
# Simplification of the equational theory

# Simplifications

$$(2\pi) = 0 = \dots \quad (\varphi_1) (\varphi_2) = (\varphi_1 + \varphi_2) \quad \text{---} [H] \text{---} [H] \text{---} = \text{---} \quad \text{---} [P(0)] \text{---} = \text{---}$$

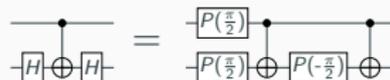
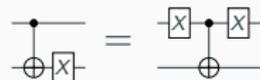


$$\text{---} [H] \text{---} = \text{---} [P(\frac{\pi}{2})] [R_X(\frac{\pi}{2})] [P(\frac{\pi}{2})] \text{---} \quad \text{---} [R_X(\alpha_1)] [P(\alpha_2)] [R_X(\alpha_3)] \text{---} = \text{---} (\beta_0) [P(\beta_1)] [R_X(\beta_2)] [P(\beta_3)] \text{---}$$

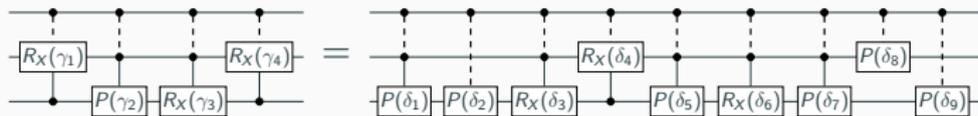
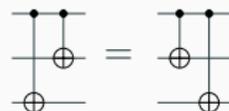
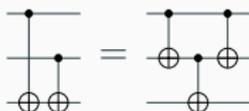


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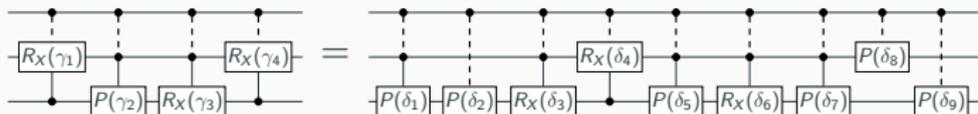
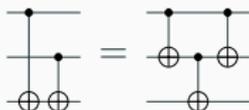
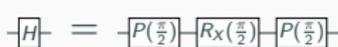
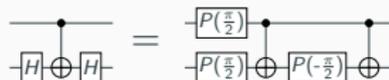
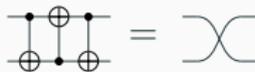
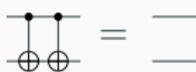


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# Simplifications

$$(2\pi) = 0 = \dots$$

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$$-|H\rangle\langle H| = -$$

$$-|P(0)\rangle\langle P(0)| = -$$

$$\text{Control} \cdot \text{P}(\varphi) = \text{P}(\varphi)$$

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$$-|R_X(\alpha_1)\rangle\langle R_X(\alpha_1)| - |P(\alpha_2)\rangle\langle P(\alpha_2)| - |R_X(\alpha_3)\rangle\langle R_X(\alpha_3)| = -|\beta_0\rangle\langle \beta_0| - |P(\beta_1)\rangle\langle P(\beta_1)| - |R_X(\beta_2)\rangle\langle R_X(\beta_2)| - |P(\beta_3)\rangle\langle P(\beta_3)|$$

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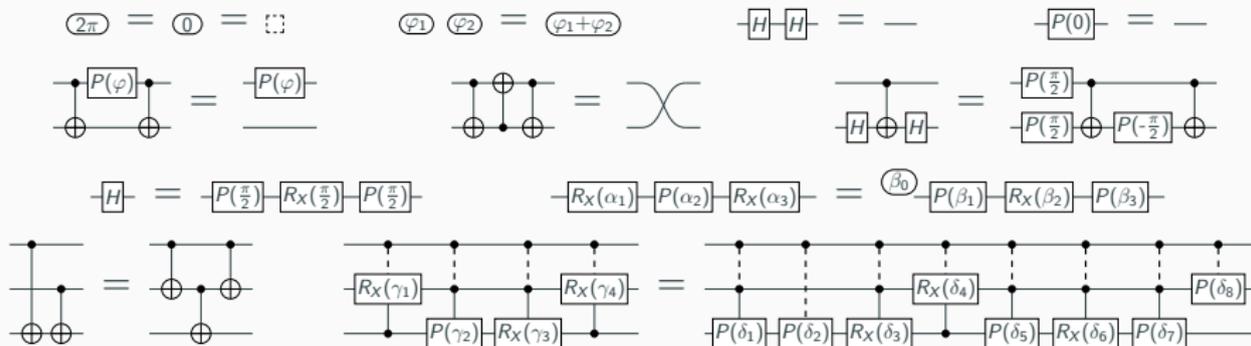
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$$R_X(\gamma_1) R_X(\gamma_4) P(\gamma_2) R_X(\gamma_3) = P(\delta_1) P(\delta_2) R_X(\delta_3) R_X(\delta_4) P(\delta_5) R_X(\delta_6) P(\delta_7) P(\delta_8) P(\delta_9)$$



# Simplified complete and sound equational theory



## Theorem (Completeness)

This equational theory is **complete**, i.e. any two equivalent circuits can be transformed into each other.

Second contribution

# Extension of the equational theory

# Extension to quantum circuits with ancillae

Quantum circuits **with ancillae** are generated by



together with



respectively denoting **qubit initialisation** and **qubit destruction**.

(The generator  $\dashv$  can only be applied to qubits in the  $|0\rangle$ -state.)

## Semantics

We extend  $\llbracket \cdot \rrbracket$  with  $\llbracket \vdash \rrbracket = |0\rangle$  and  $\llbracket \dashv \rrbracket = \langle 0|$ .

Universal for isometries

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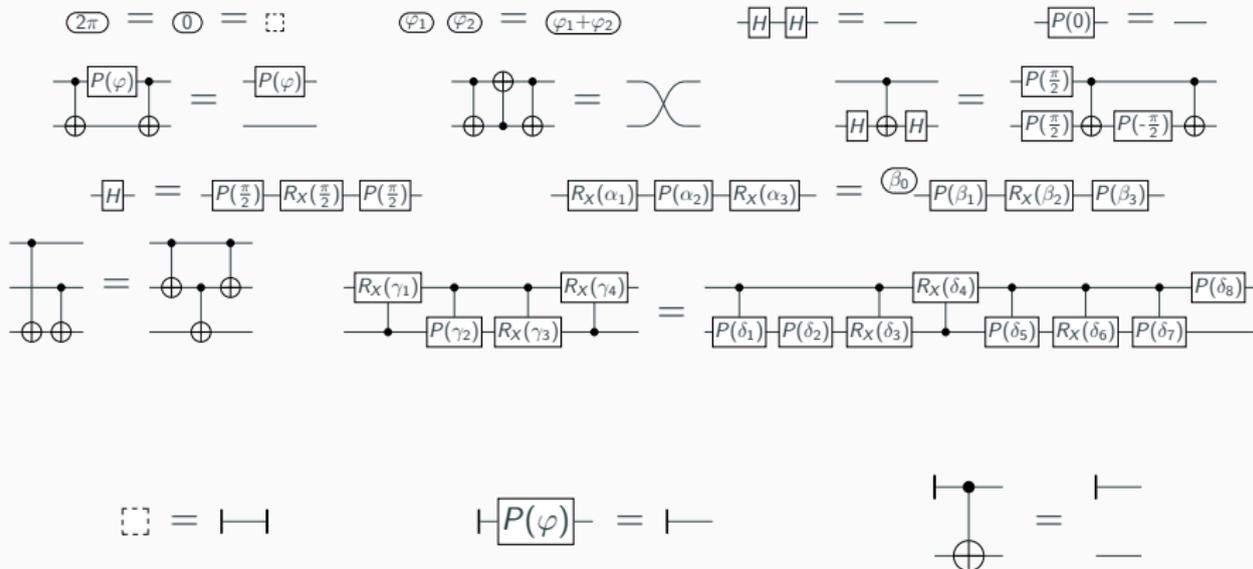
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# Equational theory for quantum circuits with ancillae

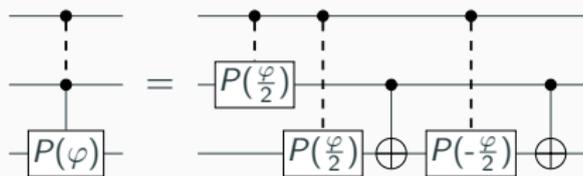


## Proposition

The big rule can be replaced by its 2-qubits version, leading to an equational theory acting on a bounded number of qubits.

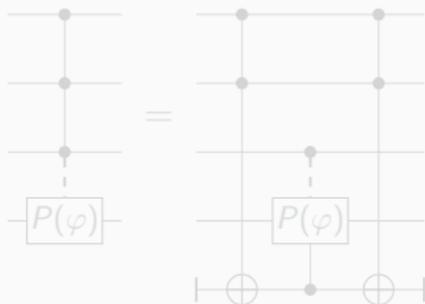
# Alternative definition of multi-controlled gates

Multi-controlled gates are defined inductively



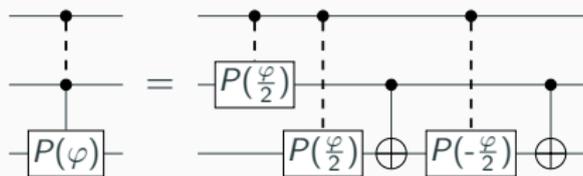
**Problem:** Cannot apply inductive hypothesis as angles are divided by 2.

**Solution:** Using ancillae, we can prove



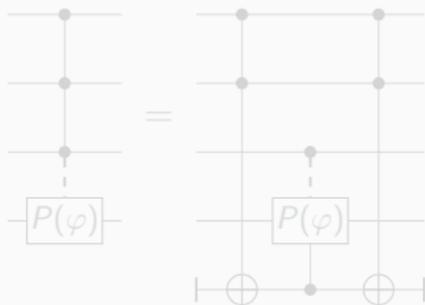
# Alternative definition of multi-controlled gates

Mutli-controlled gates are defined inductively



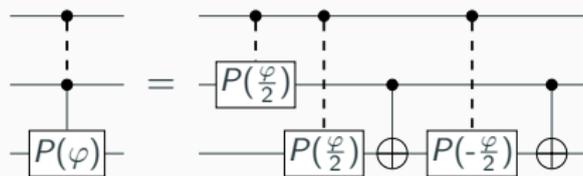
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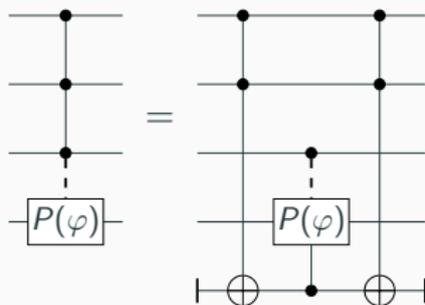
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- Simplification of the original equational theory, in particular removed two intricate rules.
- Introducing new techniques to reason on quantum circuits.
- Extension of the completeness result to circuits with ancillae.
- In these extended settings, the equational theory is made only of equations acting on at most 3 qubits.
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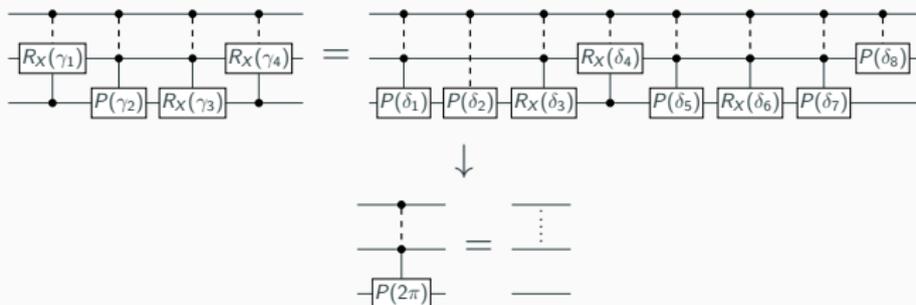
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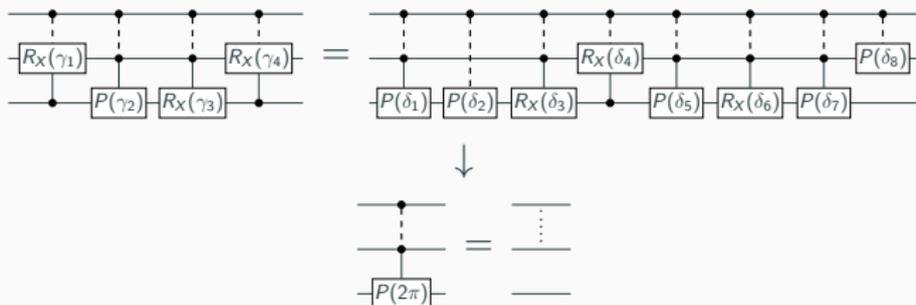


Prove the minimality of the resulting equational theory.

## Theorem (Minimality)

Each equation of the equational theory is **necessary**.

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# Thanks



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